

On Tensor Decompositions

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霍普夫代数与张量范畴 国际研讨会

International Conference on
Hopf Algebras and Tensor Categories

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- * At the categorical level, or the macro level, one usually decomposes objects into sums of simpler objects. Typical examples appear in the theories of representation rings, Brauer groups, branching laws, etc.
- * In many problems, one needs to consider explicit elements in a given object. For example, additive and multiplicative decompositions of polynomials are ubiquitous. **This talk is about decompositions of tensor elements at the micro level.**

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- * **Definition:** Let U, V, W be finite-dimensional vector spaces over some field \mathbb{K} . A 3-tensor $T \in U \otimes V \otimes W$ is said to be of **rank one**, if $T = u \otimes v \otimes w$ for some $u \in U, v \in V, w \in W$. The **rank**, denoted $R(T)$, of a tensor $T \in U \otimes V \otimes W$ is the smallest number r such that T may be expressed as a sum of r rank one tensors.

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* 3-Tensors: $\sum_{i,j,k} a_{ijk} e_i \otimes f_j \otimes g_k$ with prescribed bases

$$\iff (a_{ijk})_{l \times m \times n} \iff f(x, y, z) = \sum_{i,j,k} a_{ijk} x_i y_j z_k$$

$$\iff \Phi : U \otimes V \otimes W \rightarrow \mathbb{K}, \quad e_i \otimes f_j \otimes g_k \mapsto a_{ijk}$$

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* Call the tensor $\sum_{i_1, i_2, \dots, i_d}^n a_{i_1 i_2 \dots i_d} e_{i_1} \otimes e_{i_2} \otimes \dots \otimes e_{i_d} \in V^{\otimes d}$ symmetric, if the constants $a_{i_1 i_2 \dots i_d}$ are invariant under permuting the indices.

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- * Symmetric tensors are synonymous with homogeneous polynomials, forms, symmetric multilinear forms, projective hypersurfaces, etc.

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* For symmetric $T = \sum_{i_1, i_2, \dots, i_d}^n a_{i_1 i_2 \dots i_d} e_{i_1} \otimes e_{i_2} \otimes \dots \otimes e_{i_d}$

$\iff T = v_1^{\otimes d} + v_2^{\otimes d} + \dots + v_r^{\otimes d}$ for some $v_i \in V$

$\iff f(x_1, x_2, \dots, x_n) = \sum_{i_1, i_2, \dots, i_d}^n a_{i_1 i_2 \dots i_d} x_{i_1} x_{i_2} \dots x_{i_d} = \sum_{k=1}^r l_i^d$

where $l_i = \sum_{j=1}^n p_{ij} x_j \iff v_i = \sum_{j=1}^n p_{ij} e_j$ for all $1 \leq i \leq r$.

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* The previous expression of f is called a Waring decomposition.

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- * Remark 1: It is extremely hard, in fact NP-hard, to determine the rank of a general high-order tensor.
- * Remark 2: For symmetric tensors, one can consider the problem of symmetric tensor rank. This is another version of the Waring problem of polynomials.
- * Remark 3: A rank one tensor is essentially a monomial in the more familiar terminology of algebra. It is nontrivial to determine if a polynomial is equivalent to a monomial.

Motivating Example 1: Sums of Squares

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* A beautiful sums of squares identity:

$$(x_1^2 + x_2^2)(y_1^2 + y_2^2) = (x_1y_1 + x_2y_2)^2 + (x_1y_2 - x_2y_1)^2.$$

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- * For what values of n can there be identities like

$$(x_1^2 + \cdots + x_n^2)(y_1^2 + \cdots + y_n^2) = z_1^2 + \cdots + z_n^2$$

with $z_j \in \mathbb{K}[x_1, \dots, x_n, y_1, \dots, y_n]$?

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n-square identities exist only for $n=1, 2, 4, 8$.
- * Hurwitz's Problem (1898):
Given r and s , find minimal n such that there holds
 $(x_1^2 + \cdots + x_r^2)(y_1^2 + \cdots + y_s^2) = z_1^2 + \cdots + z_n^2$.



Motivating Example 1: Sums of Squares



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* Hopf condition and Hopf's theorem:

$$(x_1^2 + \cdots + x_r^2)(y_1^2 + \cdots + y_s^2) = z_1^2 + \cdots + z_n^2 \implies n \geq r * s$$

where $r * s$ is the Hopf-Stiefel function defined as

$$r * s := \min\{m \in \mathbb{N} \mid (x + y)^m = 0 \in \mathbb{F}_2[x, y]/(x^r, y^s)\}.$$



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* **The square identity induces a map of projective space**
 $\mathbb{P}^{r-1} \times \mathbb{P}^{s-1} \rightarrow \mathbb{P}^{n-1}$, **thus a map of the cohomology rings**
 $\mathbb{F}_2[T]/(T^n) \rightarrow \mathbb{F}_2[R]/(R^r) \otimes \mathbb{F}_2[S]/(S^s), \quad T \mapsto R \otimes 1 + 1 \otimes S.$
Here, $\mathbb{F}_2[T]/(T^n)$ is the cohomology ring of \mathbb{P}^{n-1} . Then Hopf used the Hopf structures on the cohomology rings to deduce his condition. This is one of the historical origins of Hopf algebras!



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- * On the other hand, the law of moduli of finite-dimensional real division algebras give rise to the 1, 2, 4, 8-square identities.
- * We noticed that **the octonions are essentially a commutative associative algebra in suitable braided tensor categories**. Then we could vastly generalize octonion algebras and applied them to generate square identities via partial law of moduli.
- * Furthermore, it is natural to ask the composition law of general quadratic forms, or even higher degree forms, for example

$$(x_1^d + \cdots + x_r^d)(y_1^d + \cdots + y_s^d) = z_1^d + \cdots + z_n^d.$$

Motivating Example 2: Matrix Multiplication

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- * Consider the multiplication of 2 by 2 matrices. As a bilinear map, it can be described as a 3-tensor:

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$$M = \sum_{1 \leq i, j, k \leq 2} \epsilon_{ij} \otimes \epsilon_{jk} \otimes e_{ik}.$$

- * Strassen found that $R(M) \leq 7$! This implies that, it takes seven or even less scalar multiplications to compute the product of two 2 by 2 matrices. Therefore, it takes $n^{\log_2 7}$ scalar multiplications for that of n by n matrices.

Motivating Example 2: Matrix Multiplication

- * The Strassen algorithm can be described as:

$$\begin{aligned} M = & (\epsilon_{11} + \epsilon_{22}) \otimes (\epsilon_{11} + \epsilon_{22}) \otimes (e_{11} + e_{22}) \\ & + (\epsilon_{21} + \epsilon_{22}) \otimes \epsilon_{11} \otimes (e_{21} - e_{22}) \\ & + \epsilon_{11} \otimes (\epsilon_{12} - \epsilon_{22}) \otimes (e_{12} + e_{22}) \\ & + \epsilon_{22} \otimes (\epsilon_{21} - \epsilon_{11}) \otimes (e_{11} + e_{21}) \\ & + (\epsilon_{11} + \epsilon_{12}) \otimes \epsilon_{22} \otimes (e_{12} - e_{11}) \\ & + (\epsilon_{21} - \epsilon_{11}) \otimes (\epsilon_{11} + \epsilon_{12}) \otimes e_{22} \\ & + (\epsilon_{12} - \epsilon_{22}) \otimes (\epsilon_{21} + \epsilon_{22}) \otimes e_{11}. \end{aligned}$$

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<https://doi.org/10.1038/s41586-022-05172-4>

Received: 2 October 2021

Accepted: 2 August 2022

Published online: 5 October 2022

Open access

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Improving the efficiency of algorithms for fundamental computations can have a widespread impact, as it can affect the overall speed of a large amount of computations. Matrix multiplication is one such primitive task, occurring in many systems—from neural networks to scientific computing routines. The automatic discovery of algorithms using machine learning offers the prospect of reaching beyond human intuition and outperforming the current best human-designed algorithms. However, automating the algorithm discovery procedure is intricate, as the space of possible algorithms is enormous. Here we report a deep reinforcement learning approach based on AlphaZero¹ for discovering efficient and provably correct algorithms for the multiplication of arbitrary matrices. Our agent, AlphaTensor, is trained to play a single-player game where the objective is finding tensor decompositions within a finite factor space. AlphaTensor discovered algorithms that outperform the state-of-the-art complexity for many matrix sizes. **Particularly relevant is the case of 4×4 matrices in a finite field, where AlphaTensor’s algorithm improves on Strassen’s two-level algorithm for the first time, to our knowledge, since its discovery 50 years ago².** We further showcase the flexibility of AlphaTensor through different use-cases: algorithms with state-of-the-art complexity for structured matrix multiplication and improved practical efficiency by optimizing matrix multiplication for runtime on specific hardware. Our results highlight AlphaTensor’s ability to accelerate the process of algorithmic discovery on a range of problems, and to optimize for different criteria.

We focus on the fundamental task of matrix multiplication, and use deep reinforcement learning (DRL) to search for provably correct and efficient matrix multiplication algorithms. This algorithm discovery process is particularly amenable to automation because a rich space of matrix multiplication algorithms can be formalized as low-rank decompositions of a specific three-dimensional (3D) tensor², called the matrix multiplication tensor^{3–7}. This space of algorithms contains the standard matrix multiplication algorithm and recursive algorithms such as Strassen’s², as well as the (unknown) asymptotically optimal algorithm. Although an important body of work aims at characterizing the complexity of the asymptotically optimal algorithm^{8–12}, this does not yield practical algorithms⁵. We focus here on practical matrix multiplication algorithms, which correspond to explicit low-rank decompositions of the matrix multiplication tensor. In contrast to two-dimensional matrices, for which efficient polynomial-time algorithms computing the rank have existed for over two centuries¹³, **finding low-rank decompositions of 3D tensors (and beyond) is NP-hard¹⁴** and is also hard in practice. In fact, the search space is so large that even the optimal algorithm for multiplying two 3×3 matrices is still unknown. Nevertheless, in a longstanding research effort, matrix multiplication algorithms have

been discovered by attacking this tensor decomposition problem using human search^{2,15,16}, continuous optimization^{17–19} and combinatorial search²⁰. These approaches often rely on human-designed heuristics, which are probably suboptimal. We instead use DRL to learn to recognize and generalize over patterns in tensors, and use the learned agent to predict efficient decompositions.

We formulate the matrix multiplication algorithm discovery procedure (that is, the tensor decomposition problem) as a single-player game, called **TensorGame**. At each step of TensorGame, the player selects how to combine different entries of the matrices to multiply. A score is assigned based on the number of selected operations required to reach the correct multiplication result. This is a challenging game with an enormous action space (more than 10^{12} actions for most interesting cases) that is much larger than that of traditional board games such as chess and Go (hundreds of actions). To solve TensorGame and find efficient matrix multiplication algorithms, we develop a DRL agent, AlphaTensor. AlphaTensor is built on AlphaZero^{1,21}, where a neural network is trained to guide a planning procedure searching for efficient matrix multiplication algorithms. Our framework uses a single agent to decompose matrix multiplication tensors of various sizes, yielding

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AlphaEvolve: A coding agent for scientific and algorithmic discovery

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In this white paper, we present *AlphaEvolve*, an evolutionary coding agent that substantially enhances capabilities of state-of-the-art LLMs on highly challenging tasks such as tackling open scientific problems or optimizing critical pieces of computational infrastructure. *AlphaEvolve* orchestrates an autonomous pipeline of LLMs, whose task is to improve an algorithm by making direct changes to the code. Using an evolutionary approach, continuously receiving feedback from one or more evaluators, *AlphaEvolve* iteratively improves the algorithm, potentially leading to new scientific and practical discoveries. We demonstrate the broad applicability of this approach by applying it to a number of important computational problems. When applied to optimizing critical components of large-scale computational stacks at Google, *AlphaEvolve* developed a more efficient scheduling algorithm for data centers, found a functionally equivalent simplification in the circuit design of hardware accelerators, and accelerated the training of the LLM underpinning *AlphaEvolve* itself. Furthermore, *AlphaEvolve* discovered novel, provably correct algorithms that surpass state-of-the-art solutions on a spectrum of problems in mathematics and computer science, significantly expanding the scope of prior automated discovery methods (Romera-Paredes et al., 2023). **Notably, AlphaEvolve developed a search algorithm that found a procedure to multiply two 4×4 complex-valued matrices using 48 scalar multiplications; offering the first improvement, after 56 years, over Strassen’s algorithm in this setting.** We believe *AlphaEvolve* and coding agents like it can have a significant impact in improving solutions of problems across many areas of science and computation.

1. Introduction

Discovering new high-value knowledge, such as making a novel scientific discovery or developing a commercially valuable algorithm, generally requires a prolonged process of ideation, exploration, backtracking on unpromising hypotheses, experimentation, and validation. There has been much recent interest in using large language models (LLMs) to automate significant parts of this process. Hopes of success here are driven by the breathtaking power of recent LLMs [31, 75], which can enhance their capabilities using test-time compute, and the rise of *agents* that combine language generation and action [87, 113]. These advances have improved performance across a range of established benchmarks and accelerated discovery-oriented tasks like hypothesis generation [33] and experiment design [7, 42]. However, getting LLM pipelines all the way to making entirely new scientific or practical discoveries remains challenging.

In this white paper, we present an LLM code superoptimization agent, called *AlphaEvolve*, that takes on this challenge using a combination of evolutionary computation and LLM-based code generation. *AlphaEvolve* focuses on the broad spectrum of scientific and engineering

¹See Acknowledgments and Author information section. ^{*}Equal contributions.

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- * Problem 2: How to decompose a tensor as a direct product, i.e., a multiplicative decomposition?
- * Problem 3: Determine if a tensor is of rank one, or if a homogeneous polynomial is a product of linear forms.

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- * A tensor $T \in U \otimes V \otimes W$ is called a direct sum, if there are $T_i \in U_i \otimes V_i \otimes W_i$ ($i = 1, 2$) such that $T = T_1 + T_2$ in the tensor space $(U_1 \oplus U_2) \otimes (V_1 \oplus V_2) \otimes (W_1 \oplus W_2) = U \otimes V \otimes W$.

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- * In terms of polynomials, this is equivalent to writing $f(X, Y, Z) = g(X_1, Y_1, Z_1) + h(X_2, Y_2, Z_2)$ where $X = X_1 \sqcup X_2$, etc. In other words, separate variables of polynomials.

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Then

$$R(T) = R(T_1) + R(T_2).$$

- * The conjecture is disproved by Shitov (Acta Math. 2019). It seems that the conjecture fails only in very strange situations. It is of interest to consider the conditions for “=” holds.

Direct Product

- * **Definition:** A multivariate function $f(x_1, x_2, \dots, x_n)$ is called a direct product if, after a linear change of variables, it can be written as a product of two or more functions in disjoint sets of variables

$$f(x) \stackrel{x=Py}{=} g_1(y_1, \dots, y_a) \cdot g_2(y_{a+1}, \dots, y_n), \quad 0 < a < n.$$

Call $f(x_1, x_2, \dots, x_n)$ completely factorable, if it is a product of n univariate functions.

Direct Sum vs Direct Product

- * Lemma. (1) $f(x_1, x_2, \dots, x_n)$ is a direct sum if and only if \exp^f is a direct product. (2) $g(x_1, x_2, \dots, x_n)$ is a direct product if and only if $\log g$ is a direct sum.

Center and Decompositions

* **Definition:** The center of an n -variate function $f(x_1, x_2, \dots, x_n)$ is defined as

$$Z(f) := \{X \in \mathbb{K}^{n \times n} \mid (HX)^T = HX\},$$

where $H = \left(\frac{\partial^2 f}{\partial x_i \partial x_j} \right)_{1 \leq i, j \leq n}$ is the hessian matrix of the function $f(x_1, x_2, \dots, x_n)$.

Center and Decompositions

- * Proposition: (1) The center of a multivariate function is a special Jordan algebra. (2) Direct sum decompositions of f are in bijection with complete sets of orthogonal idempotent of $Z(f)$. (3) Direct product decompositions of f are in bijection with complete sets of orthogonal idempotents of $Z(\log f)$.

Toy Example 1

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* If $f = x^3 + y^3$, then its center $Z(f) \cong \mathbb{K}^2$.

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- * By dimension counting:
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- * Cardano formula revisited:
$$a_0x^3 + 3a_1x^2 + 3a_2x + a_3 = (\alpha_1x + \beta_1)^3 + (\alpha_2x + \beta_2)^3$$

Toy Example 2

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$$f(x_1, x_2, x_3) = x_1^3 + x_2^3 + x_3^3 - 3x_1x_2x_3 .$$

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- * The center of the function $\log f$ is

$$Z(\log f) = \mathbb{C} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} + \mathbb{C} \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{pmatrix} + \mathbb{C} \begin{pmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}.$$

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- * The polynomial f is factorized as
$$f(x_1, x_2, x_3) = (x_1 + x_2 + x_3)(x_1 + \omega x_2 + \omega^2 x_3)(x_1 + \omega^2 x_2 + \omega x_3)$$
where $\omega = \frac{-1 + \sqrt{3}i}{2}$ is a primitive cubic root of unity.

The Chow Variety

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- * Problem: determine whether a higher degree form is a product of linear forms.

The Chow Variety

The Chow Variety

- * **Theorem:** Suppose $f \in \mathbb{K}[x_1, x_2, \dots, x_n]_d$. Then $f = l_1^{d_1} l_2^{d_2} \dots l_n^{d_n}$ ($d_i \geq 1, \forall i$) with linearly independent l_i 's if and only if $Z(\log f) \cong \mathbb{K}^n$.

The Chow Variety

- * **Theorem:** Suppose $f \in \mathbb{K}[x_1, x_2, \dots, x_n]_d$. Then $f = l_1^{d_1} l_2^{d_2} \dots l_n^{d_n}$ ($d_i \geq 1, \forall i$) with linearly independent l_i 's if and only if $Z(\log f) \cong \mathbb{K}^n$.
- * **Remark:** Our algorithm of direct product gives a simple criterion and an efficient algorithm to completely factorize a multivariate polynomial.

Toy Example 3

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- * By calculation, it is easy to observe that $\partial_x \partial_y (\partial_x - \partial_y)$ annihilates the previous quartic. Moreover, each linear factor of the annihilating differential polynomial corresponds to a term of the Waring decomposition!

Hopf Pairing vs The Apolarity Method

Hopf Pairing vs The Apolarity Method

- * Denote $\mathcal{R} = \mathbb{K}[x_1, x_2, \dots, x_n]$ and $\mathcal{D} = \mathbb{K}[\partial_1, \partial_2, \dots, \partial_n]$. Then \mathcal{D} acts naturally as differentiation on \mathcal{R} .

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- * \mathcal{R} and \mathcal{D} are Hopf algebras, and they are mutually graded dual of each other. The previous action comes from Hopf algebraic structures and it provides a nice Hopf pairing.

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- * \mathcal{R} and \mathcal{D} are Hopf algebras, and they are mutually graded dual of each other. The previous action comes from Hopf algebraic structures and it provides a nice Hopf pairing.
- * Rota highly advocated the Hopf algebraic approach to the apolarity method and even the whole classical algebraic invariant theory!

Hopf Pairing vs The Apolarity Method

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- * For $f \in \mathcal{R}_d$, let $f^\perp = \{g \in \mathcal{D} \mid g \circ f = 0\}$. Then f^\perp is a **Gorenstein** ideal and \mathcal{D}/f^\perp is a Gorenstein algebra. A well known result of **Macaulay** says that: **there is a bijection between forms and graded Artinian Gorenstein algebras.**

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- * If $g \in \mathcal{D}_d$, then $g^\perp = \{f \in \mathcal{R} \mid g \circ f = 0\} \subset \mathcal{R}$ is a coideal. **Determining g^\perp amounts to solving a constant coefficient PDE!** For example, consider the Laplacian $\Delta = \partial_{x_1}^2 + \dots + \partial_{x_n}^2$, then Δ^\perp consists of the well-known harmonic polynomials.

Hopf Pairing vs The Apolarity Method

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* The Apolarity Lemma: Suppose $f \in \mathbb{C}[x_1, \dots, x_n]$. Then

$$f = \sum_{i=1}^r \lambda_i (\alpha_{i1}x_1 + \dots + \alpha_{in}x_n)^d \iff \bigcap_{i=1}^r I_{[\alpha_{i1} : \dots : \alpha_{in}]} \subset f^\perp.$$

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* For binary forms, we provide an elementary treatment. For forms with more variables, it seems Hopf pairings might be a promising approach.

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- * The problem of Tensor decompositions arises naturally from arithmetic, algebra, geometry, multilinear algebra, data processing, and various other practical areas.
- * There remain many challenging problems. Methods from Hopf algebras and tensor categories, representation theory, invariant theory, algebraic geometry, and even machine learning, artificial intelligence, etc. play significant roles for further studying.

Thank You!

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Questions or Comments?